

ADT Specification for a "List".

List

Operation : ADD :

\* Pre Cond : A List should be exist. ↖ 12, 15, 17

Input : The item to be added. — 25

Process : The item will be added to the end of the list.

Output : Success of addition / None.

Post Cond : List with an added item. — 12, 15, 17, 25

As ADT Specifications we should understand,

- \* ① How Data is stored.
- \* ② Operations.

Arrays

	•			
	•			
x = 0	5	[1][1]	[1][1]	
1	7	[1][2]	[2][1]	
2	12	[1][3]	[3][1]	
3	2	[1][4]	[4][1]	
4	8	[1][5]	[5][1]	
	-	[2][1]	[1][2]	
	-	[2][2]	[2][2]	
	-	[2][3]	[3][2]	
	•			

Row major ordering.      Column major ordering.

Student Index No	S <sub>1</sub> M <sub>1</sub>	S <sub>2</sub> M	S <sub>3</sub> M	...
1	25	75	98	58
2	75	25	72	73
3	50	65	83	28
4	55	* 85	44	22
5	100	100	86	31

2 dimension

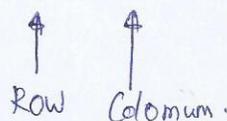
A[i][j]

= x + (i-1) \* Size of a row + (j-1) \* Size of an element.

ex: If I want \*access\* value,  
= x + (4-1) \* Size of a row + (3-1) \* Size of an element

row major ordering.

A [1][2]



3 ways we can store data.

\* 2D Array of String

\* Array of Records.

\* Different arrays for different records.

### ADT - Linked List

2 structure implementing a LinkedList.

① Record / struct Node {

int : data.

\* Node: next.

};

② Record / struct LinkedList {

\* Node head

};

ex: LinkedList

LIST\_CREATE () {

LinkedList l = new LinkedList ();

l → head = NULL

return l.

};

This Linked List L  
can have 2 scenarios:

① head → null

② head → [1] → [3]

After adding x.

① head → [x]

② head → [x] → [1] → [3]

Void LIST\_INSERT ( LinkedList L, int x ) {  
method      method name      Parameters.

Node newNode = newNode ();

newNode → data = x

newNode → next = head;

L → head = newNode. }

Above Process for,

①st Scenario.

head → NULL

newNode → [ ]

newNode → [x]

newNode → [x]

②nd Scenario.

head → [1] → [3]

newNode → [ ]

newNode → [x]

newNode → [x] → [1] → [3]

```

Boolean LIST_SEARCH (LinkedList L, int x) {
    Node temp = L → head
    while (temp != NULL) {
        if (temp → data == x)
            return TRUE
        temp = temp → next;
    }
    return FALSE
}

```

```

Void LIST_DELETE (LinkedList L, int x) {

```

If the 1st element we have to delete

```

    if (L → head == NULL)
        return.
    if (L → head → data == x) {
        Node temp = L → head
        L → head = L → head → next
        delete (temp)
        return.
    }

```

If its in the middle.

```

    Node Previous = head.
    while (prev → next != NULL) {
        if (prev → next → data == x) {
            Node temp = prev → next
            prev → next = temp → next
            delete (temp)
            return.
        }

```

```

        prev = prev → next

```

```

    }

```

```

}

```

## Doubly linked list

Slide

```
List CREATE() {  
    DoublyLinkedList l = new LinkedList();  
    l → head = NULL  
    return l.  
}
```

```
Void LIST_INSERT (DoublyLinkedList L, int x) {  
    Node newNode = newNode(x)  
    newNode → data = x  
    newNode → next = head;  
    newNode → Prev = NULL  
    head → Prev = newNode  
    L → head = newNode.  
}
```

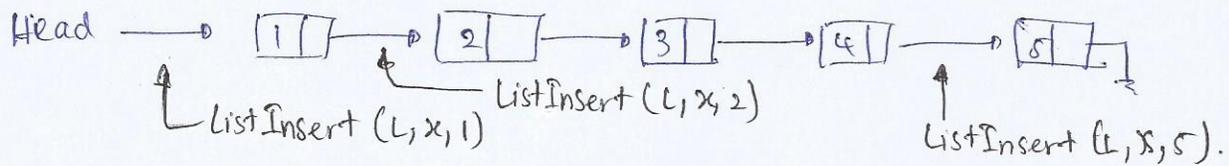
Searching is Same.

In Deletion,

```
⋮  
If L → head → data == x {  
    L → head → next → prev = NULL.  
    L → head = L → head → next,  
    ⋮  
    Node temp = prev → next  
    temp → next → prev = temp → prev.  
    prev → next = temp → next.  
    ⋮  
    ⋮  
    ⋮
```

## List Insertion for $i$ th.

Pg - (3)



List\_Insert (LinkedList L, int x, int location)

else Node newNode = newNode();      if (loc == 1)  
List\_Insert(L, x)

newNode → data = x

int i = 2;

curr = head

while (i < loc && curr → next != NULL)

curr = curr → next;

NewN → next = curr → next

curr → next = NewN

## Queue

struct Node {

data

next

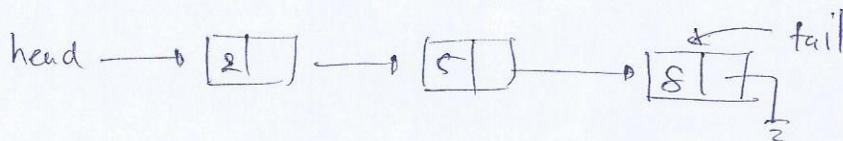
};

struct Queue {

\* Node head

\* Node tail

};



## Stacks

struct Stack {

Node top

};